

**Question 1:**

[1 point]

Give five distinguishable strings for the language  $L = \{0,1\}^*\{00\}\{0,1\}$  and show why they are pairwise distinguishable.

**Question 2:**

[2 points]

- Give a *deterministic* finite automaton  $M_1$  recognizing the language  $L_1$  over the alphabet  $\Sigma = \{a,b\}$  containing all strings with no two consecutive equal alphabet symbols.
- Construct a *deterministic* finite automaton  $M_2$  recognizing the language  $L_2$  denoted by the regular expression  $(a+ab)^*bb^*$ .
- Give a *deterministic* finite automaton  $M_3$  recognizing the complement of the above language  $L_2$ .

**Question 3:**

[1,5 points]

- Construct a *nondeterministic* finite automaton  $M$  recognizing the language  $L(G)$  generated by the regular grammar  $G$  with the following productions

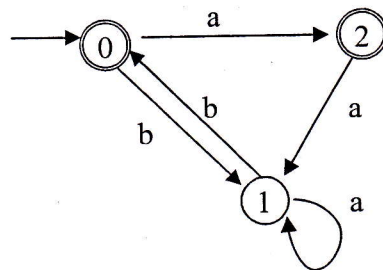
$$S \rightarrow bA \quad A \rightarrow aA \mid aS \mid aB \mid \Lambda \quad B \rightarrow bS \mid \Lambda$$

- Use the *powerset construction* to construct a *deterministic* finite automaton  $N$  recognizing the same language as the above *nondeterministic* automaton  $M$ .
- Use the automaton  $N$  above to construct a *regular grammar* generating the same language accepted by  $N$ . It may be helpful to rename the states of the automaton  $N$ .

**Question 4:**

[1 point]

Use the *state elimination* method of Brzozowski and McCluskey to construct a regular expression for the language recognized by the following finite automaton



**Question 5:**

[1 point]

Use the pumping lemma to show that the language  $L = \{a^n b^n a^m \mid n \geq 0, m \geq n\}$  is not regular.

**Question 6:**

[2 points]

- Find a context-free grammar  $G_1$  generating the language  $L_1 = \{a^n b^n \mid n > 0\}$ .
- Find a context-free grammar  $G_2$  generating the language  $L_2 = \{b^n \mid n > 0\}$ .
- Use the above two grammar to give a context free grammar for the language  $L_3 = L_1 \cdot L_2$ .
- Find a context-free grammar  $G$  generating the language  $L_4 = L_3^*$ .

**Question 7:**

[1.5 points]

- Draw a pushdown automaton  $M$  recognizing the language  $L = \{a^n b^m \mid m > n > 0\}$  using as alphabet symbols only  $A$  and  $Z_0$  (the initial stack symbol).
- Use the above pushdown automaton  $M$  to construct a new pushdown automaton  $M_e$  accepting the above language  $L$  by *empty stack* (thus without accepting states).

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The final score is given by the sum of the points obtained.